


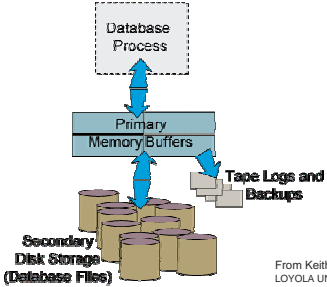
Storage and File Structure II

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Big Picture




From Keith Van Rhein's slide, LOYOLA UNIVERSITY CHICAGO

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File Organization

- The database is stored as a collection of *files*. Each file is a sequence of *records*. A record is a sequence of fields.
- Two cases:
 - Fixed length record
 - Variable length record



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Fixed-Length Records



record 0	A-102	Perryridge	400
record 1	A-305	Round Hill	350
record 2	A-215	Mianus	700
record 3	A-101	Downtown	500
record 4	A-222	Redwood	700
record 5	A-201	Perryridge	900
record 6	A-217	Brighton	750
record 7	A-110	Downtown	600
record 8	A-218	Perryridge	700

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Free Lists



header				
record 0	A-102	Perryridge	400	
record 1				
record 2	A-215	Mianus	700	
record 3	A-101	Downtown	500	
record 4				
record 5	A-201	Perryridge	900	
record 6				
record 7	A-110	Downtown	600	
record 8	A-218	Perryridge	700	

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Variable-Length Records



- Variable-length records arise in database systems in several ways:
 - Storage of multiple record types in a file.
 - Record types that allow variable lengths for one or more fields.
 - Record types that allow repeating fields (used in some older data models).

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Byte String Representation



0	Perryridge	A-102	400	A-201	900	⊥
1	Round Hill	A-305	350	⊥		
2	Mianus	A-215	700	⊥		
3	Downtown	A-101	500	A-110	600	⊥
4	Redwood	A-222	700	⊥		
5	Brighton	A-217	750	⊥		



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Dealing with Variable-Length Record

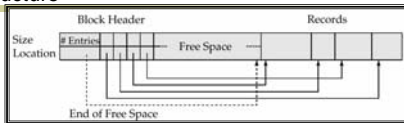


- By introducing pointers
- Stuff empty fields

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Variable-Length Records: Slotted Page Structure



- Slotted page header contains:
 - number of record entries
 - end of free space in the block
 - location and size of each record

Compare and contrast Slotted Page Structure with Byte String Representation

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Variable-Length Records - Fixed-length Representation



0	Perryridge	A-102	400	A-201	900	A-218	700
1	Round Hill	A-305	350	⊥	⊥	⊥	⊥
2	Mianus	A-215	700	⊥	⊥	⊥	⊥
3	Downtown	A-101	500	A-110	600	⊥	⊥
4	Redwood	A-222	700	⊥	⊥	⊥	⊥
5	Brighton	A-217	750	⊥	⊥	⊥	⊥

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Pointer Method



0	Perryridge	A-102	400	
1	Round Hill	A-305	350	
2	Mianus	A-215	700	
3	Downtown	A-101	500	
4	Redwood	A-222	700	
5		A-201	900	
6	Brighton	A-217	750	
7		A-110	600	
8		A-218	700	

- Pointer method
 - A variable-length record is represented by a list of fixed-length records, chained together via pointers.
 - Can be used even if the maximum record length is not known

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Pointer Method (Cont.)



- Waste space
- Solution is to allow two kinds of block in file:
 - Anchor block – contains the first records of chain
 - Overflow block – contains records other than those that are the first records of chains.

anchor block	Perryridge	A-102	400	
	Round Hill	A-305	350	
	Mianus	A-215	700	
	Downtown	A-101	500	
	Redwood	A-222	700	
	Brighton	A-217	750	
overflow block		A-201	900	
		A-218	700	
		A-110	600	

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Organization of Records in Files



- **Heap** – no order
- **Sequential** – sequential order based on search key
- **Hashing** – a hash function computed on some attribute of each record; the result specifies in which block of the file the record should be placed
- **Clustering file organization** – records of several different relations can be stored in the same file
 - Motivation: store related records on the same block to minimize I/O

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Sequential File Organization



Search key

A-217	Brighton	750	→
A-101	Downtown	500	→
A-110	Downtown	600	→
A-215	Mianus	700	→
A-102	Perryridge	400	→
A-201	Perryridge	900	→
A-218	Perryridge	700	→
A-222	Redwood	700	→
A-305	Round Hill	350	→

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Sequential File Organization (Cont.)



- Deletion – use pointer chains
- Insertion – may need buffer
- Need to reorganize the file from time to time to restore sequential order

A-217	Brighton	750	→
A-101	Downtown	500	→
A-110	Downtown	600	→
A-215	Mianus	700	→
A-102	Perryridge	400	→
A-201	Perryridge	900	→
A-218	Perryridge	700	→
A-222	Redwood	700	→
A-305	Round Hill	350	→
A-888	North Town	800	→

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Clustering File Organization

	Hayes	Main	Brooklyn
Customer	Hayes	A-102	
	Hayes	A-220	
	Hayes	A-503	
Account	Turner	Putnam	Stamford
	Turner	A-305	

Advantages and disadvantages?

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Clusters in Oracle

- **Cluster**
 CREATE CLUSTER personnel (department NUMBER(4));
 CREATE CLUSTER personnel_hash (department NUMBER(4)) HASH IS
 department HASHKEYS 200;
- **Cluster Keys**
 CREATE INDEX idx_personnel ON CLUSTER personnel;
- **Adding Tables to a Cluster**
 CREATE TABLE dept(
 department number(4),
 name char(60),
 address char(40))
 CLUSTER personnel (department);
 CREATE TABLE faculty (
 name char(60),
 address char(40),
 department number(4))
 CLUSTER personnel (department);

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Oracle's Data Blocks, Extents and Segments

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Oracle Block Structure

Overhead includes header information, a row directory and table directory.

The block will have a certain amount of free space

Data records fill the used portion of the block

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Storage option in Oracle

- create table student (
 - student_id number,
 - name char(60),
 - address char(40)
 storage (
 - initial 50K
 - next 50K
 - maxextents 10
 - pctincrease 25);

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Data Dictionary Storage

Data dictionary (also called system catalog) stores metadata

- Information about relations
- User and accounting information, including passwords
- Statistical and descriptive data
 - number of tuples in each relation
- Physical file organization information
- Information about indices

Oracle Demo

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Data Dictionary Storage (Cont.)



- Catalog structure: can use either
 - specialized data structures designed for efficient access
 - a set of relations, with existing system features used to ensure efficient access

The latter alternative is usually preferred

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